Characterizing possible typical asymptotic behaviours of cellular automata

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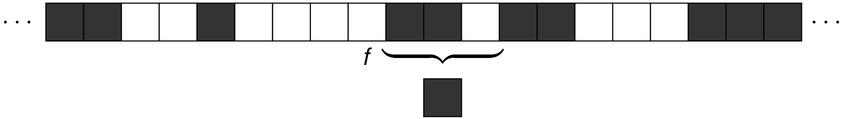
Definitions

- $\mathcal{A}, \mathcal{B}$ finite **alphabets**;
- $\mathcal{A}^*$ the (finite) **words**;
- $\mathcal{A}^{\mathbb{Z}}$ the **configurations**;
- σ the **shift action** $\sigma(a)_i = a_{i-1}$;

A **cellular automaton** is an action $F : \mathcal{A}^{\mathbb{Z}} \to \mathcal{A}^{\mathbb{Z}}$ defined by a **local rule** $f : \mathcal{A}^{\mathbb{U}} \to \mathcal{A}$ on some neighbourhood $\mathbb{U}$.

For $\mathcal{A} = \{\blacksquare, \square\}$ and $\mathbb{U} = \{-1, 0, 1\}$:



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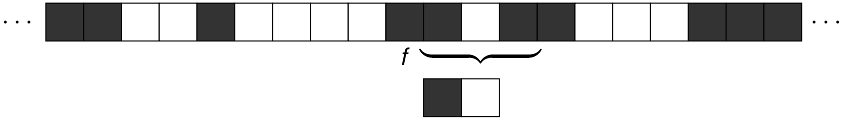

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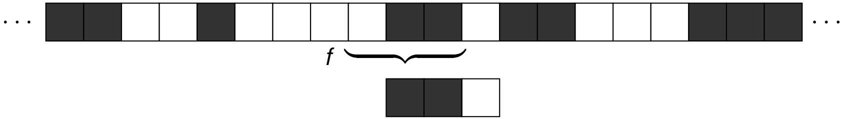

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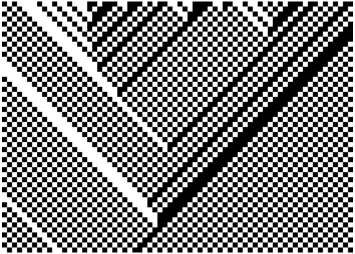
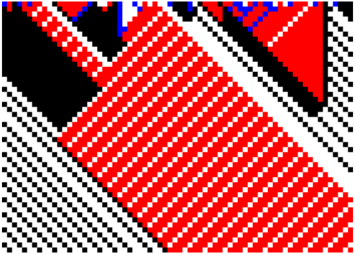
Simulations and typical asymptotic behaviour

Traffic automaton



Captive automaton



3-state cyclic automaton

Additive automaton

Measure space

$\mathcal{M}_\sigma(\mathcal{A}^\mathbb{Z})$ the **σ-invariant probability measures** on $\mathcal{A}^\mathbb{Z}$.

$\mu([u])$ the probability that a word $u \in \mathcal{A}^*$ appears, for $\mu \in \mathcal{M}_\sigma(\mathcal{A}^\mathbb{Z})$.

Examples

Bernoulli (i.i.d) measures Let $(\lambda_a)_{a\in\mathcal{A}}$ such that $\sum \lambda_a = 1$.

$$\boxed{\forall u \in \mathcal{A}^*, \mu([u]) = \prod_{i=0}^{|u|-1} \lambda_{u_i}.}$$

Measures supported by a periodic orbit For a finite word w,

$$\boxed{\widehat{\delta_w} = \frac{1}{|w|} \sum_{i=0}^{|w|-1} \delta_{\sigma^i({}^\infty w^\infty)}.}$$

Markov measures with finite memory.

Action of an automaton on an initial measure

- F extends to an action $F_* : \mathcal{M}_\sigma(\mathcal{A}^{\mathbb{Z}}) \to \mathcal{M}_\sigma(\mathcal{A}^{\mathbb{Z}})$:

$$\boxed{F_*\mu(U) = \mu(F^{-1}U)}$$

for any borelian U.

- For an initial measure μ, $F_*^t\mu$ describes the repartition at time t;
- Typical asymptotic behaviour is well described by the limit(s) of $(F_*^t\mu)_{t\in\mathbb{N}}$ in the **weak-* topology**:

$$\boxed{F_*^t\mu \underset{t\to\infty}{\longrightarrow} \nu \quad \Leftrightarrow \quad \forall u \in \mathcal{A}^*, F_*^t\mu([u]) \to \nu([u]).}$$

Examples of asymptotic behaviour

Examples of asymptotic behaviour

Proposition

Let μ be the uniform Bernoulli measure on $\{0, 1, 2\}$ and F the 3-state cyclic automaton.

$$\boxed{F_*^t\mu \to \frac{1}{3}\widehat{\delta_0} + \frac{1}{3}\widehat{\delta_1} + \frac{1}{3}\widehat{\delta_2}.}$$

Main question

Question

Which measures ν are reachable as the limit of the sequence $(F_*^t\mu)_{t\in\mathbb{N}}$ for some cellular automaton F and initial measure μ?

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Answer

All (take $F = Id$ and $\mu = \nu$).

Main question

Better question

Which measures ν are reachable as the limit of the sequence $(F_*^t\mu)_{t\in\mathbb{N}}$ for some cellular automaton F and **simple** initial measure μ (e.g. the uniform Bernoulli measure)?

In a sense, this would correspond to the "physically relevant" measure for F.

Section 2

Necessary conditions: computability obstructions

Topological obstructions

Topological obstruction

The accumulation points of $(F_*^t\mu)_{t\in\mathbb{N}}$ form a nonempty and **compact** set.

Measures and computability

A probability measure $\mu \in \mathcal{M}_\sigma(\mathcal{A}^{\mathbb{Z}})$ is:

computable if $\boxed{u \to \mu([u]) \text{ is computable}}$,

i.e. if there exists $f : \mathcal{A}^* \times \mathbb{N} \to \mathbb{Q}$ computable such that

$$\boxed{|\mu([u]) - f(u,n)| < 2^{-n}}.$$

($\Leftrightarrow$ can be **simulated** by a probabilistic Turing machine)

Examples of computable measures

- Any measure supported by a periodic orbit;
- Any Bernoulli or Markov measure with computable parameters.

Measures and computability

A probability measure $\mu \in \mathcal{M}_\sigma(\mathcal{A}^{\mathbb{Z}})$ is:

semi-computable ($\emptyset'$-computable) if there exists a computable function $f : \mathcal{A}^* \times \mathbb{N} \to \mathbb{Q}$ such that

$$\boxed{|\mu([u]) - f(u,n)| \underset{n\to\infty}{\longrightarrow} 0.}$$

($\Leftrightarrow$ **limit** of a computable sequence of measures)

Examples of computable measures

- Any measure supported by a periodic orbit;
- Any Bernoulli or Markov measure with computable parameters.

Computability obstruction

Action of an automaton on a computable measure

- If μ is computable, then $F_*^t\mu$ is **computable**;
- If μ is computable, and $F_*^t\mu \underset{t\to\infty}{\longrightarrow} \nu$,
  then ν is **semi-computable**.

Section 3

Sufficient conditions: construction of limit measures

State of the art

Motto:

"The only obstruction is the computability obstruction"

Theorem [Hochman, Meyerovitch 10]

Possible entropies for multidimensional subshifts of finite type are exactly the reals aproximable from above.

Theorem [Boyer, Poupet, Theyssier 06], [Boyer, Delacourt, Sablik 10]

The language of words u satisfying

$$F_*^t\mu([u]) \not\to 0$$

can be Σ_3-**complete** for any nondegenerate Bernoulli measure μ.

Main result

Action of an automaton on a computable measure

If μ is computable, and $F_*^t\mu \underset{t\to\infty}{\longrightarrow} \nu$, then ν is **semi-computable**.

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Action of an automaton on a computable measure

If μ is computable, and $F_*^t\mu \underset{t\to\infty}{\longrightarrow} \nu$, then ν is **semi-computable**.

Theorem

Let ν be a **semi-computable** measure. There exists:

- an alpabet $\mathcal{B} \supset \mathcal{A}$
- a cellular automaton $F : \mathcal{B} \to \mathcal{B}$

such that, for any **ergodic** and **full-support** measure $\mu \in \mathcal{M}_\sigma(\mathcal{B}^\mathbb{Z})$,

$$\boxed{F_*^t\mu \underset{t\to\infty}{\longrightarrow} \nu}$$

Approximation by periodic orbits

Proposition

Measures supported by periodic orbits are dense in $\mathcal{M}_\sigma(\mathcal{A}^{\mathbb{Z}})$.

Example: Uniform Bernoulli measure

$w_0 = 01$
$w_1 = 0011$
$w_2 = 00010111$
$w_3 = 0000110100101111$

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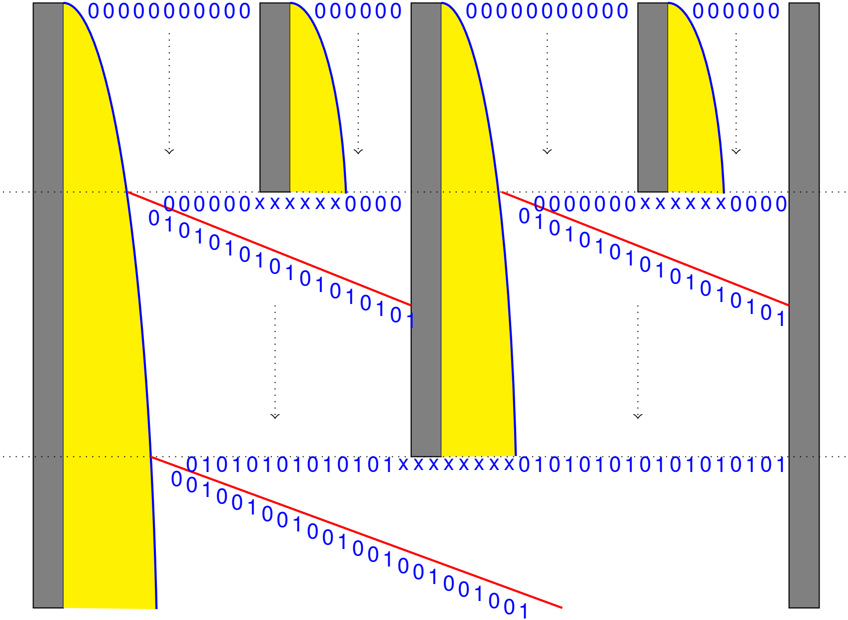
Proposition

If $\nu \in \mathcal{M}_\sigma(\mathcal{A}^{\mathbb{Z}})$ is semi-computable, there is a **computable** sequence of words $(w_n)_{n\in\mathbb{N}}$ such that $\widehat{\delta_{w_n}} \to \nu$.

Our construction will compute each w_n and approach the measure $\widehat{\delta_{w_n}}$ by writing concatenated copies of w_n on all the configuration.



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Section 4

Extensions and related results

Extensions and related results

Questions

1. Sets of accumulation points? Yes, with a computability condition on compact sets

Compact sets and computability

Consider the following distance function:

$$d_{\mathcal{M}}(\mu_1, \mu_2) = \sum_{n=0}^{\infty} \frac{1}{2^n} \max_{u \in \mathcal{A}^n} |\mu_1([u]) - \mu_2([u])|$$

Then the **computability of a compact set** $\mathcal{V}$ can be defined in the following way.

$\mathcal{V}$ computable if $d_{\mathcal{V}} : \mathcal{M}_\sigma(\mathcal{A}^{\mathbb{Z}}) \mapsto \mathbb{R}$ is **computable**, that is:

$$\exists f : \mathcal{A}^* \times \mathbb{N} \to \mathbb{Q} \text{ computable}, |d_{\mathcal{V}}(\widehat{\delta_w}) - f(w, n)| \leq \frac{1}{2^n}$$

and $\exists b : \mathbb{N} \mapsto \mathbb{Q}$ computable,

$$d_{\mathcal{M}}(\mu_1, \mu_2) < b(m) \Rightarrow |d_{\mathcal{V}}(\mu_1) - d_{\mathcal{V}}(\mu_2)| \leq \frac{1}{2^m}$$

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Then the **computability of a compact set** $\mathcal{V}$ can be defined in the following way.

$\mathcal{V}$ $\emptyset'$-lower-semi-computable if $\boxed{\boldsymbol{d_{\mathcal{V}} = \liminf d_i}}$, where d_i are elements in $\mathcal{M}_\sigma(\mathcal{A}^{\mathbb{Z}}) \mapsto \mathbb{R}$, and:

$\exists f : \mathbb{N} \times \mathcal{A}^* \times \mathbb{N} \mapsto \mathbb{Q}$ computable,

$$|d_i(\widehat{\delta_w}) - f(i, w, n)| \leq \frac{1}{2^n} \textbf{ (sequential computability)}$$

and $\exists b : \mathbb{N} \mapsto \mathbb{Q}$ computable,

$$d_{\mathcal{M}}(\mu_1, \mu_2) < b(m) \Rightarrow |d_i(\mu_1) - d_i(\mu_2)| \leq \frac{1}{2^m} \textbf{ (effective uniform equicontinuity).}$$

Computability obstructions, again

Action of an automaton on a computable measure:

- If μ is computable, then $F_*^t\mu$ is **computable**;
- If μ is computable and the accumulation points of $(F_*^t\mu)_{t\in\mathbb{N}}$ are $\mathcal{V}$, then $\mathcal{V}$ is nonempty, compact and **$\emptyset'$-lower-semi-computable**.

Intuitively, $d_{\mathcal{V}} = \liminf d_{\mathcal{M}}(F_*^t\mu, .)$.

Theorem

Let $\mathcal{V}$ be a nonempty, compact, **connected**, **$\emptyset'$-lower-semi-computable** set of measures.
Then there exists an automaton $F : \mathcal{A} \to \mathcal{A}$ such that, for any measure $\mu \in \mathcal{M}_\sigma(\mathcal{A}^{\mathbb{Z}})$ **σ-mixing** and **full-support**,

$$\boxed{\text{The set of accumulation points of } (F_*^t\mu)_{t\in\mathbb{N}} \text{ is } \mathcal{V}.}$$

Extensions and related results

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Extensions and related results

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1. Sets of accumulation points? Yes, with a computability condition on compact sets
2. Implementation of the construction? No (but for good reasons)

Implementation

- Non-trivial Turing machine satisfying space constraints;
- Large number of states;
  (for $|\mathcal{B}| = 2$, at least 2244 times more than the corresponding Turing machine)
- Speed of convergence $O\left(\frac{1}{\log t}\right)$ in the best case.

Extensions and related results

Questions

1. Sets of accumulation points? Yes, with a computability condition on compact sets
2. Implementation of the construction? No (but for good reasons)
3. No auxiliary states? Yes, if the target measure is not full-support

Theorem

Let ν be a **non full-support**, **semi-computable** measure.
Then there exists an automaton $F : \mathcal{A} \to \mathcal{A}$ such that, for any measure $\mu \in \mathcal{M}_\sigma(\mathcal{A}^\mathbb{Z})$ **σ-mixing** and **full-support**,

$$\boxed{F_*^t \mu \underset{t\to\infty}{\longrightarrow} \nu.}$$

Extensions and related results

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3. No auxiliary states? Yes, if the target measure is not full-support

Idea: use forbidden words to encode auxiliary states.

Remark

If $F_*^t\mu \to \nu$ where ν is a full support measure, then F is a **surjective** automaton and **the uniform Bernoulli measure is invariant**.

Extensions and related results

Questions

1. Sets of accumulation points? Yes, with a computability condition on compact sets
2. Implementation of the construction? No (but for good reasons)
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4. Cesaro mean convergence? Yes

Extensions and related results

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3. No auxiliary states? Yes, if the target measure is not full-support
4. Cesaro mean convergence? Yes
5. Characterization of the support? In progress

Conjecture

K is a Σ_n-computable compact set of measures

$$\Leftrightarrow$$

$\overline{\bigcup_{\nu \in K} supp(\nu)}$ is a Σ_{n+1}-computable compact set of configurations.

Extensions and related results

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6. Properties of the limit measure? Mostly undecidable

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6. Properties of the limit measure? Mostly undecidable
7. Using the initial measure as an argument or an oracle? Some simple cases

Computation in the space of measures

Let us consider the operator

$$\boxed{\mu \mapsto \text{accumulation points of } (F_*^t \mu)_{t\in\mathbb{R}}}$$

The previous construction gave us operators that were essentially **constant** (on a large domain).

Question

Which operators $\mathcal{M}_\sigma(\mathcal{A}^\mathbb{Z}) \to \mathcal{M}_\sigma(\mathcal{A}^\mathbb{Z})$ (ou $\mathcal{M}_\sigma(\mathcal{A}^\mathbb{Z}) \to \mathcal{P}(\mathcal{M}_\sigma(\mathcal{A}^\mathbb{Z}))$) can be realized in this way?

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Theorem

Let $\nu : \mathbb{R} \to \mathcal{M}_\sigma(\mathcal{A}^\mathbb{Z})$ be a **semi-computable** operator. There is:

- an alphabet $\mathcal{B} \supset \mathcal{A}$,
- an automaton $F : \mathcal{B}^\mathbb{Z} \to \mathcal{B}^\mathbb{Z}$

such that, for any **full-support** and **exponentially σ-mixing** measure μ,

$$\boxed{F_*^t\mu \underset{t\to\infty}{\longrightarrow} \nu\left(\mu\left(\boxed{\mathrm{I}}\right)\right)}.$$

Some examples

Let $M \subset \mathcal{M}_\sigma(\mathcal{B}^{\mathbb{Z}})$ be the set of **full-support**, **exponentially σ-mixing** measures.

Example 1: Density classification

There exists an automaton $F : \mathcal{B}^{\mathbb{Z}} \to \mathcal{B}^{\mathbb{Z}}$ realizing the operator:

$$M \to \mathcal{M}_\sigma(\{0,1\}^{\mathbb{Z}})$$
$$\mu \mapsto \begin{cases} \widehat{\delta_0} & \text{if } \mu(\boxed{1}) < \frac{1}{2} \\ \widehat{\delta_1} & \text{otherwise.} \end{cases}$$

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\end{aligned}
$$

Example 2: A simple oracle

There exists an automaton $F : \mathcal{B}^{\mathbb{Z}} \to \mathcal{B}^{\mathbb{Z}}$ realizing the operator:

$$
\begin{aligned}
M &\to \mathcal{M}_\sigma(\{0,1\}^{\mathbb{Z}}) \\
\mu &\mapsto Ber(\mu(\boxed{1}))
\end{aligned}
$$

Implementation of a simple case

Fibonacci word

Consider the morphism :

$$\begin{array}{rccc} \varphi : & \mathcal{A}^* & \to & \mathcal{A}^* \\ & 0 & \mapsto & 01 \\ & 1 & \mapsto & 0 \end{array}$$

Then the sequence $\varphi^n(0)$ converges to an infinite word called **Fibonacci word**:

$$\varphi^\infty(0) = 0100101001001010010101\ldots$$

and it is **uniquely ergodic**.